The 10[th] Romanian Master of Mathematics Competition

Day 1 — Solutions

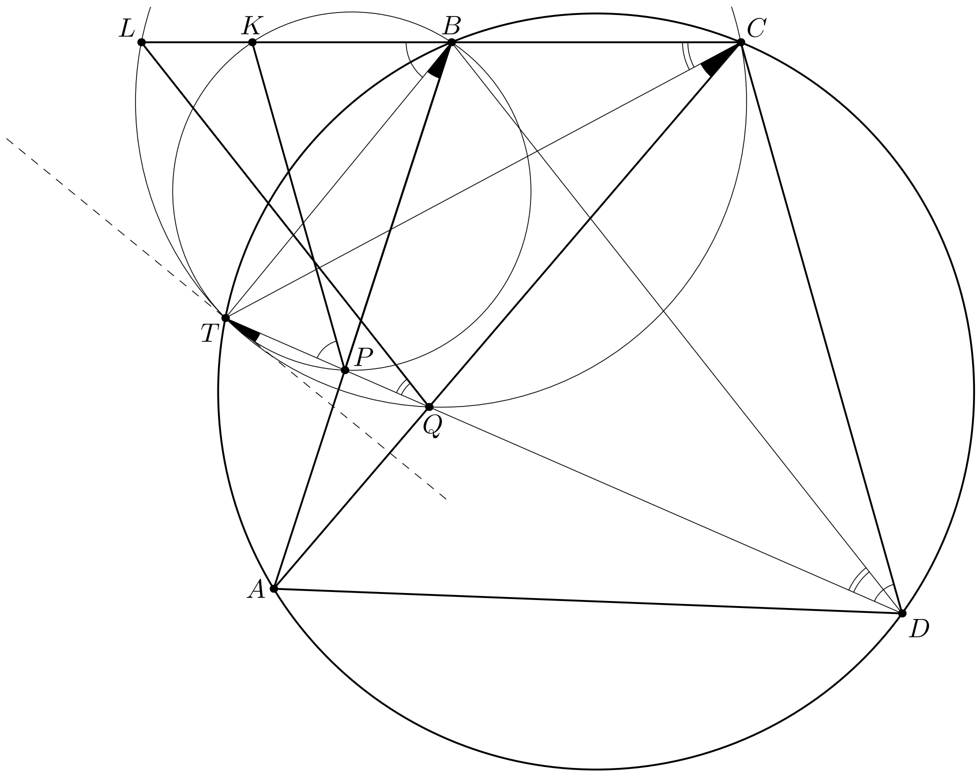
Problem 1. Let $ABCD$ be a cyclic quadrangle and let P be a point on the side AB. The diagonal AC crosses the segment DP at Q. The parallel through P to CD crosses the extension of the side BC beyond B at K, and the parallel through Q to BD crosses the extension of the side BC beyond B at L. Prove that the circumcircles of the triangles BKP and CLQ are tangent.

Aleksandr Kuznetsov, Russia

Solution. We show that the circles BKP and CLQ are tangent at the point T where the line DP crosses the circle $ABCD$ again.

Since $BCDT$ is cyclic, we have $\angle KBT = \angle CDT$. Since $KP \parallel CD$, we get $\angle CDT = \angle KPT$. Thus, $\angle KBT = \angle CDT = \angle KPT$, which shows that T lies on the circle BKP. Similarly, the equalities $\angle LCT = \angle BDT = \angle LQT$ show that T also lies on the circle CLQ.

It remains to prove that these circles are indeed tangent at T. This follows from the fact that the chords TP and TQ in the circles $BKTP$ and $CLTQ$, respectively, both lie along the same line and subtend equal angles $\angle TBP = \angle TBA = \angle TCA = \angle TCQ$.



Remarks. The point T may alternatively be defined as the Miquel point of (any four of) the five lines AB, BC, AC, KP, and LQ.

Of course, the result still holds if P is chosen on the line AB, and the other points lie on the corresponding lines rather than segments/rays. The current formulation was chosen in order to avoid case distinction based on the possible configurations of points.

Problem 2. Determine whether there exist non-constant polynomials $P(x)$ and $Q(x)$ with real coefficients satisfying

$$P(x)^{10} + P(x)^9 = Q(x)^{21} + Q(x)^{20}. \qquad (*)$$

Ilya Bogdanov, Russia

Solution 1. The answer is in the negative. Comparing the degrees of both sides in $(*)$ we get $\deg P = 21n$ and $\deg Q = 10n$ for some positive integer n. Take the derivative of $(*)$ to obtain

$$P'P^8(10P + 9) = Q'Q^{19}(21Q + 20). \qquad (**)$$

Since $\gcd(10P+9, P) = \gcd(10P+9, P+1) = 1$, it follows that $\gcd(10P+9, P^9(P+1)) = 1$, so $\gcd(10P+9, Q) = 1$, by $(*)$. Thus $(**)$ yields $10P+9 \mid Q'(21Q+20)$, which is impossible since $0 < \deg(Q'(21Q+20)) = 20n-1 < 21n = \deg(10P+9)$. A contradiction.

Remark. A similar argument shows that there are no non-constant solutions of $P^m + P^{m-1} = Q^k + Q^{k-1}$, where k and m are positive integers with $k \geq 2m$. A critical case is $k = 2m$; but in this case there exist more routine ways of solving the problem. Thus, we decided to choose $k = 2m+1$.

Solution 2. Letting r and s be integers such that $r \geq 2$ and $s \geq 2r$, we show that if $P^r + P^{r-1} = Q^s + Q^{s-1}$, then Q is constant.

Let $m = \deg P$ and $n = \deg Q$. A degree inspection in the given relation shows that $m \geq 2n$.

We will prove that $P(P+1)$ has at least $m+1$ distinct complex roots. Assuming this for the moment, notice that Q takes on one of the values 0 or -1 at each of those roots. Since $m+1 \geq 2n+1$, it follows that Q takes on one of the values 0 and -1 at more than n distinct points, so Q must be constant.

Finally, we prove that $P(P+1)$ has at least $m+1$ distinct complex roots. This can be done either by referring to the Mason–Stothers theorem or directly, in terms of multiplicities of the roots in question.

Since P and $P+1$ are relatively prime, the Mason–Stothers theorem implies that the number of distinct roots of $P(P+1)$ is greater than m, hence at least $m+1$.

For a direct proof, let $z_1, \ldots, z_t$ be the distinct complex roots of $P(P+1)$, and let z_k have multiplicity α_k, $k = 1, \ldots, t$. Since P and $P+1$ have no roots in common, and $P' = (P+1)'$, it follows that P' has a root of multiplicity $\alpha_k - 1$ at z_k. Consequently, $m - 1 = \deg P' \geq \sum_{k=1}^{t}(\alpha_k - 1) = \sum_{k=1}^{t} \alpha_k - t = 2m - t$; that is, $t \geq m+1$. This completes the prof.

Remark. The Mason–Stothers theorem (in a particular case over the complex field) claims that, given coprime complex polynomials $P(x)$, $Q(x)$, and $R(x)$, not all constant, such that $P(x) + Q(x) = R(x)$, the total number of their complex roots (**not** regarding multiplicities) is at least $\max\{\deg P, \deg Q, \deg R\} + 1$. This theorem was a part of motivation for the famous *abc*-conjecture.

Problem 3. Ann and Bob play a game on an infinite checkered plane making moves in turn; Ann makes the first move. A move consists in orienting any unit grid-segment that has not been oriented before. If at some stage some oriented segments form an oriented cycle, Bob wins. Does Bob have a strategy that guarantees him to win?

Maxim Didin, Russia

Solution. The answer is in the negative: Ann has a strategy allowing her to prevent Bob's victory.

We say that two unit grid-segments form a *low-left corner* (or *LL-corner*) if they share an endpoint which is the lowest point of one and the leftmost point of the other. An *up-right corner* (or *UR-corner*) is defined similarly. The common endpoint of two unit grid-segments at a corner is the *joint* of that corner.

Fix a vertical line on the grid and call it the *midline*; the unit grid-segments along the midline are called *middle segments*. The unit grid-segments lying to the left/right of the midline are called *left/right segments*. Partition all left segments into LL-corners, and all right segments into UR-corners.

We now describe Ann's strategy. Her first move consists in orienting some middle segment arbitrarily. Assume that at some stage, Bob orients some segment s. If s is a middle segment, Ann orients any free middle segment arbitrarily. Otherwise, s forms a corner in the partition with some other segment t. Then Ann orients t so that the joint of the corner is either the source of both arrows, or the target of both. Notice that after any move of Ann's, each corner in the partition is either completely oriented or completely not oriented. This means that Ann can always make a required move.

Assume that Bob wins at some stage, i.e., an oriented cycle C occurs. Let X be the lowest of the leftmost points of C, and let Y be the topmost of the rightmost points of C. If X lies (strictly) to the left of the midline, then X is the joint of some corner whose segments are both oriented. But, according to Ann's strategy, they are oriented so that they cannot occur in a cycle — a contradiction. Otherwise, Y lies to the right of the midline, and a similar argument applies. Thus, Bob will never win, as desired.

Remarks. (1) There are several variations of the argument in the solution above. For instance, instead of the midline, Ann may choose any infinite in both directions down going polyline along the grid (i.e., consisting of steps to the right and steps-down alone). Alternatively, she may split the plane into four quadrants, use their borders as "trash bin" (as the midline was used in the solution above), partition all segments in the upper-right quadrant into UR-corners, all segments in the lower-right quadrant into LR-corners, and so on.

(2) The problem becomes easier if Bob makes the first move. In this case, his opponent just partitions the whole grid into LL-corners. In particular, one may change the problem to say that the first player to achieve an oriented cycle wins (in this case, the result is a draw).

On the other hand, this version is closer to known problems. In particular, the following problem is known:

> *Ann and Bob play the game on an infinite checkered plane making moves in turn (Ann makes the first move). A move consists in painting any unit grid segment that has not been painted before (Ann paints in blue, Bob paints in red). If a player creates a cycle of her/his color, (s)he wins. Does any of the players have a winning strategy?*

Again, the solution is pairing strategy with corners of a fixed orientation (with a little twist for Ann's strategy — in this problem, it is clear that Ann has better chances).

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Day 2 — Solutions

Problem 4. Let a, b, c, d be positive integers such that $ad \neq bc$ and $\gcd(a,b,c,d)=1$. Prove that, as n runs through the positive integers, the values $\gcd(an+b, cn+d)$ may achieve form the set of all positive divisors of some integer.

Raul Alcantara, Peru

Solution 1. We extend the problem statement by allowing a and c take non-negative integer values, and allowing b and d to take arbitrary integer values. (As usual, the greatest common divisor of two integers is non-negative.) Without loss of generality, we assume $0 \leq a \leq c$. Let $S(a,b,c,d) = \{\gcd(an+b, cn+d) \colon n \in \mathbb{Z}_{>0}\}$.

Now we induct on a. We first deal with the inductive step, leaving the base case $a=0$ to the end of the solution. So, assume that $a>0$; we intend to find a 4-tuple (a', b', c', d') satisfying the requirements of the extended problem, such that $S(a', b', c', d') = S(a,b,c,d)$ and $0 \leq a' < a$, which will allow us to apply the induction hypothesis.

The construction of this 4-tuple is provided by the step of the Euclidean algorithm. Write $c = aq + r$, where q and r are both integers and $0 \leq r < a$. Then for every n we have

$$\gcd(an+b, cn+d) = \gcd\big(an+b, q(an+b) + rn + d - qb\big) = \gcd\big(an+b, rn + (d-qb)\big),$$

so a natural intention is to define $a' = r$, $b' = d - qb$, $c' = a$, and $d' = b$ (which are already shown to satisfy $S(a', b', c', d') = S(a,b,c,d)$). The check of the problem requirements is straightforward: indeed,

$$a'd' - b'c' = (c - qa)b - (d - qb)a = -(ad - bc) \neq 0$$

and

$$\gcd(a', b', c', d') = \gcd(c - qa, b - qd, a, b) = \gcd(c, d, a, b) = 1.$$

Thus the step is verified.

It remains to deal with the base case $a = 0$, i.e., to examine the set $S(0, b, c, d)$ with $bc \neq 0$ and $\gcd(b, c, d) = 1$. Let b' be the integer obtained from b by ignoring all primes b and c share (none of them divides $cn + d$ for any integer n, otherwise $\gcd(b, c, d) > 1$). We thus get $\gcd(b', c) = 1$ and $S(0, b', c, d) = S(0, b, c, d)$.

Finally, it is easily seen that $S(0, b', c, d)$ is the set of all positive divisors of b'. Each member of $S(0, b', c, d)$ is clearly a divisor of b'. Conversely, if δ is a positive divisor of b', then $cn + d \equiv \delta \pmod{b'}$ for some n, since b' and c are coprime, so δ is indeed a member of $S(0, b', c, d)$.

Solution 2. (*Alexander Betts*) For positive integers s and t and prime p, we will denote by $\gcd_p(s,t)$ the greatest common p-power divisor of s and t.

Claim 1. For any positive integer n, $\gcd(an+b, cn+d) \mid ad - bc$.

Proof. This is clear from the identity

$$a(cn+d) - c(an+b) = ad - bc. \qquad (\dagger)$$

Claim 2. The set of values taken by $\gcd(an+b, cn+d)$ is exactly the set of values taken by the product

$$\prod_{p \mid ad-bc} \gcd_p(an_p + b, cn_p + d)$$

as the $(n_p)_{p \mid ad-bc}$ each range over positive integers.

Proof. From the identity

$$\gcd(an+b, cn+d) = \prod_{p \mid ad-bc} \gcd_p(an+b, cn+d),$$

it is clear that every value taken by $\gcd(an+b, cn+d)$ is also a value taken by the product (with all $n_p = n$). Conversely, it suffices to show that, given any positive integers $(n_p)_{p\mid ad-bc}$, there is a positive integer n such that $\gcd_p(an+b, cn+d) = \gcd_p(an_p+b, cn_p+d)$ for each $p \mid ad-bc$. This can be achieved by requiring that n be congruent to n_p modulo a sufficiently large[1] power of p (using the Chinese Remainder Theorem).

Using Claim 2, it suffices to determine the sets of values taken by $\gcd_p(an+b, cn+d)$ as n ranges over all positive integers. There are two cases.

Claim 3. If $p \mid a, c$, then $\gcd_p(an+b, cn+d) = 1$ for all n.

Proof. If $p \mid an+b, cn+d$, then we would have $p \mid a, b, c, d$, which is not the case.

Claim 4. If $p \nmid a$ or $p \nmid c$, then the values taken by $\gcd_p(an+b, cn+d)$ are exactly the p-power divisors of $ad-bc$.

Proof. Assume without loss of generality that $p \nmid a$. Then from identity (†) we have $\gcd_p(an+b, cn+d) = \gcd_p(an+b, ad-bc)$. But since $p \nmid a$, the arithmetic progression $an+b$ takes all possible values modulo the highest p-power divisor of $ad-bc$, and in particular the values taken by $\gcd_p(an+b, ad-bc)$ are exactly the p-power divisors of $ad-bc$.

Conclusion. Using claims 2, 3 and 4, we see that the set of values taken by $\gcd(an+b, cn+d)$ is equal to the set of products of p-power divisors of $ad-bc$, where we only consider those primes p not dividing $\gcd(a, c)$. Thus the set of values of $\gcd(an+b, cn+d)$ is equal to the set of divisors of the largest factor of $ad-bc$ coprime to $\gcd(a, c)$.

Remarks. (1) If $S(a, b, c, d)$ is the set of all positive divisors of some integer, then necessarily $ad \neq bc$ and $\gcd(a, b, c, d) = 1$: finiteness of $S(a, b, c, d)$ forces the former, and membership of 1 forces the latter.

(2) One may modify the problem statement according to the first paragraph of the solution. However, it seems that in this case one needs to include a clarification of the agreement on gcd being necessarily non-negative.

[1]For example, $n \equiv n_p$ modulo the largest p-power divisor of $ad-bc$.

Problem 5. Let n be a positive integer and fix $2n$ distinct points on a circumference. Split these points into n pairs and join the points in each pair by an arrow (i.e., an oriented line segment). The resulting configuration is *good* if no two arrows cross, and there are no arrows $\overrightarrow{AB}$ and $\overrightarrow{CD}$ such that $ABCD$ is a convex quadrangle *oriented clockwise*. Determine the number of good configurations.

Fedor Petrov, Russia

Solution 1. The required number is $\binom{2n}{n}$. To prove this, trace the circumference counterclockwise to label the points a_1, a_2, ..., a_{2n}.

Let $\mathcal{C}$ be any good configuration and let $O(\mathcal{C})$ be the set of all points *from* which arrows emerge. We claim that every n-element subset S of $\{a_1, \ldots, a_{2n}\}$ is an O-image of a unique good configuration; clearly, this provides the answer.

To prove the claim induct on n. The base case $n = 1$ is clear. For the induction step, consider any n-element subset S of $\{a_1, \ldots, a_{2n}\}$, and assume that $S = O(\mathcal{C})$ for some good configuration $\mathcal{C}$. Take any index k such that $a_k \in S$ and $a_{k+1} \notin S$ (assume throughout that indices are cyclic modulo $2n$, i.e., $a_{2n+1} = a_1$ etc.).

If the arrow from a_k points to some a_ℓ, $k+1 < \ell$ $(< 2n+k)$, then the arrow pointing to a_{k+1} emerges from some a_m, m in the range $k+2$ through $\ell - 1$, since these two arrows do not cross. Then the arrows $a_k \to a_\ell$ and $a_m \to a_{k+1}$ form a prohibited quadrangle. Hence, $\mathcal{C}$ contains an arrow $a_k \to a_{k+1}$.

On the other hand, if any configuration $\mathcal{C}$ contains the arrow $a_k \to a_{k+1}$, then this arrow cannot cross other arrows, neither can it occur in prohibited quadrangles.

Thus, removing the points a_k, a_{k+1} from $\{a_1, \ldots, a_{2n}\}$ and the point a_k from S, we may apply the induction hypothesis to find a unique good configuration $\mathcal{C}'$ on $2n-2$ points compatible with the new set of sources (i.e., points from which arrows emerge). Adjunction of the arrow $a_k \to a_{k+1}$ to $\mathcal{C}'$ yields a unique good configuration on $2n$ points, as required.

Solution 2. Use the counterclockwise labelling a_1, a_2, ..., a_{2n} in the solution above.

Letting D_n be the number of good configurations on $2n$ points, we establish a recurrence relation for the D_n. To this end, let $C_n = \frac{(2n)!}{n!(n+1)!}$ the nth Catalan number; it is well-known that C_n is the number of ways to connect $2n$ given points on the circumference by n pairwise disjoint chords.

Since no two arrows cross, in any good configuration the vertex a_1 is connected to some a_{2k}. Fix k in the range 1 through n and count the number of good configurations containing the arrow $a_1 \to a_{2k}$. Let $\mathcal{C}$ be any such configuration.

In $\mathcal{C}$, the vertices a_2, ..., a_{2k-1} are paired off with one other, each arrow pointing from the smaller to the larger index, for otherwise it would form a prohibited quadrangle with $a_1 \to a_{2k}$. Consequently, there are C_{k-1} ways of drawing such arrows between a_2, ..., a_{2k-1}.

On the other hand, the arrows between a_{2k+1}, ..., a_{2n} also form a good configuration, which can be chosen in D_{n-k} ways. Finally, it is easily seen that any configuration of the first kind and any configuration of the second kind combine together to yield an overall good configuration.

Thus the number of good configurations containing the arrow $a_1 \to a_{2k}$ is $C_{k-1}D_{n-k}$. Clearly, this is also the number of good configurations containing the arrow $a_{2(n-k+1)} \to a_1$, so

$$D_n = 2\sum_{k=1}^{n} C_{k-1}D_{n-k}. \qquad (*)$$

To find an explicit formula for D_n, let $d(x) = \sum_{n=0}^{\infty} D_n x^n$ and let $c(x) = \sum_{n=0}^{\infty} C_n x^n = \frac{1-\sqrt{1-4x}}{2x}$ be the generating functions of the D_n and the C_n, respectively. Since $D_0 = 1$, relation $(*)$

yields $d(x) = 2xc(x)d(x) + 1$, so

$$d(x) = \frac{1}{1 - 2xc(x)} = (1-4x)^{-1/2} = \sum_{n\geq 0} \left(-\frac{1}{2}\right)\left(-\frac{3}{2}\right)\cdots\left(-\frac{2n-1}{2}\right)\frac{(-4x)^n}{n!}$$
$$= \sum_{n\geq 0} \frac{2^n(2n-1)!!}{n!}x^n = \sum_{n\geq 0} \binom{2n}{n}x^n.$$

Consequently, $D_n = \binom{2n}{n}$.

Solution 3. Let $C_n = \frac{1}{n+1}\binom{2n}{n}$ denote the nth Catalan number and recall that there are exactly C_n ways to join $2n$ distinct points on a circumference by n pairwise disjoint chords. Such a configuration of chords will be referred to as a *Catalan n-configuration*. An orientation of the chords in a Catalan configuration $\mathcal{C}$ making it into a good configuration (in the sense defined in the statement of the problem) will be referred to as a *good* orientation for $\mathcal{C}$.

We show by induction on n that there are exactly $n+1$ good orientations for any Catalan n-configuration, so there are exactly $(n+1)C_n = \binom{2n}{n}$ good configurations on $2n$ points. The base case $n = 1$ is clear.

For the induction step, let $n > 1$, let $\mathcal{C}$ be a Catalan n-configuration, and let ab be a chord of minimal length in $\mathcal{C}$. By minimality, the endpoints of the other chords in $\mathcal{C}$ all lie on the major arc ab of the circumference.

Label the $2n$ endpoints 1, 2, …, $2n$ counterclockwise so that $\{a, b\} = \{1, 2\}$, and notice that the good orientations for $\mathcal{C}$ fall into two disjoint classes: Those containing the arrow $1 \to 2$, and those containing the opposite arrow.

Since the arrow $1 \to 2$ cannot be involved in a prohibited quadrangle, the induction hypothesis applies to the Catalan $(n-1)$-configuration formed by the other chords to show that the first class contains exactly n good orientations.

Finally, the second class consists of a single orientation, namely, $2 \to 1$, every other arrow emerging from the smaller endpoint of the respective chord; a routine verification shows that this is indeed a good orientation. This completes the induction step and ends the proof.

Remark. Combining the arguments from Solutions 1 and 3 one gets a way (though not the easiest) to compute the Catalan number C_n.

Solution 4, sketch. (*Sang-il Oum*) As in the previous solution, we intend to count the number of good orientations of a Catalan n-configuration.

For each such configuration, we consider its *dual graph* T whose vertices are finite regions bounded by chords and the circle, and an edge connects two regions sharing a boundary segment. This graph T is a plane tree with n edges and $n+1$ vertices.

There is a canonical bijection between orientations of chords and orientations of edges of T in such a way that each chord crosses an edge of T from the right to the left of the arrow on that edge. A good orientation of chords corresponds to an orientation of the tree containing no two edges oriented towards each other. Such an orientation is defined uniquely by its *source vertex*, i.e., the unique vertex having no in-arrows.

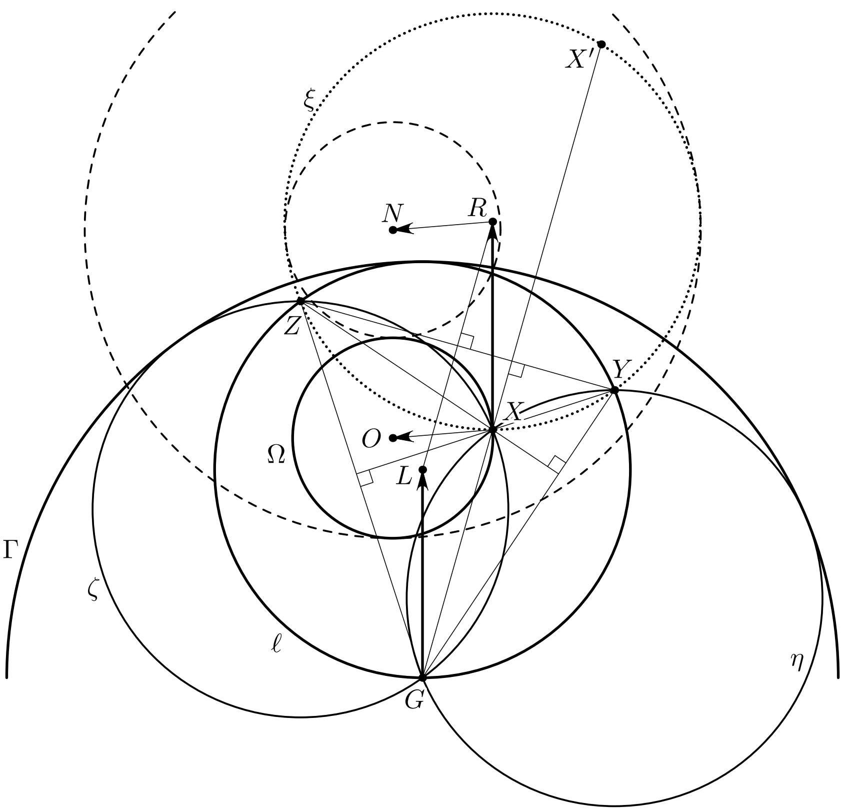
Therefore, for each tree T on $n+1$ vertices, there are exactly $n+1$ ways to orient it so that the source vertex is unique — one for each choice of the source. Thus, the answer is obtained in the same way as above.

Problem 6. Fix a circle Γ, a line ℓ tangent to Γ, and another circle Ω disjoint from ℓ such that Γ and Ω lie on opposite sides of ℓ. The tangents to Γ from a variable point X on Ω cross ℓ at Y and Z. Prove that, as X traces Ω, the circle XYZ is tangent to two fixed circles.

Russia, Ivan Frolov

Solution. Assume Γ of unit radius and invert with respect to Γ. No reference will be made to the original configuration, so images will be denoted by the same letters. Letting Γ be centred at G, notice that inversion in Γ maps tangents to Γ to circles of unit diameter through G (hence internally tangent to Γ). Under inversion, the statement reads as follows:

> *Fix a circle Γ of unit radius centred at G, a circle ℓ of unit diameter through G, and a circle Ω inside ℓ disjoint from ℓ. The circles η and ζ of unit diameter, through G and a variable point X on Ω, cross ℓ again at Y and Z, respectively. Prove that, as X traces Ω, the circle XYZ is tangent to two fixed circles.*



Since η and ζ are the reflections of the circumcircle ℓ of the triangle GYZ in its sidelines GY and GZ, respectively, they pass through the orthocentre of this triangle. And since η and ζ cross again at X, the latter is the orthocentre of the triangle GYZ. Hence the circle ξ through X, Y, Z is the reflection of ℓ in the line YZ; in particular, ξ is also of unit diameter.

Let O and L be the centres of Ω and ℓ, respectively, and let R be the (variable) centre of ξ. Let GX cross ξ again at X'; then G and X' are reflections of one another in the line YZ, so $GLRX'$ is an isosceles trapezoid. Then $LR \parallel GX$ and $\angle(LG, GX) = \angle(GX', X'R) = \angle(RX, XG)$, i.e., $LG \parallel RX$; this means that $GLRX$ is a parallelogram, so $\overrightarrow{XR} = \overrightarrow{GL}$ is constant.

Finally, consider the fixed point N defined by $\overrightarrow{ON} = \overrightarrow{GL}$. Then $XRNO$ is a parallelogram, so the distance $RN = OX$ is constant. Consequently, ξ is tangent to the fixed circles centred at N of radii $|1/2 - OX|$ and $1/2 + OX$.

One last check is needed to show that the inverse images of the two obtained circles are indeed circles and not lines. This might happen if one of them contained G; we show that this is impossible. Indeed, since Ω lies inside ℓ, we have $OL < 1/2 - OX$, so

$$NG = \left|\overrightarrow{GL} + \overrightarrow{LO} + \overrightarrow{ON}\right| = \left|2\overrightarrow{GL} + \overrightarrow{LO}\right| \geq 2\left|\overrightarrow{GL}\right| - \left|\overrightarrow{LO}\right| > 1 - (1/2 - OX) = 1/2 + OX;$$

this shows that G is necessarily outside the obtained circles.

Remarks. (1) The last check could be omitted, if we allowed in the problem statement to regard a line as a particular case of a circle. On the other hand, the Problem Selection Committee suggests not to punish students who have not performed this check.

(2) Notice that the required fixed circles are also tangent to Ω.